

e-graphs
for
next-gen PL tools

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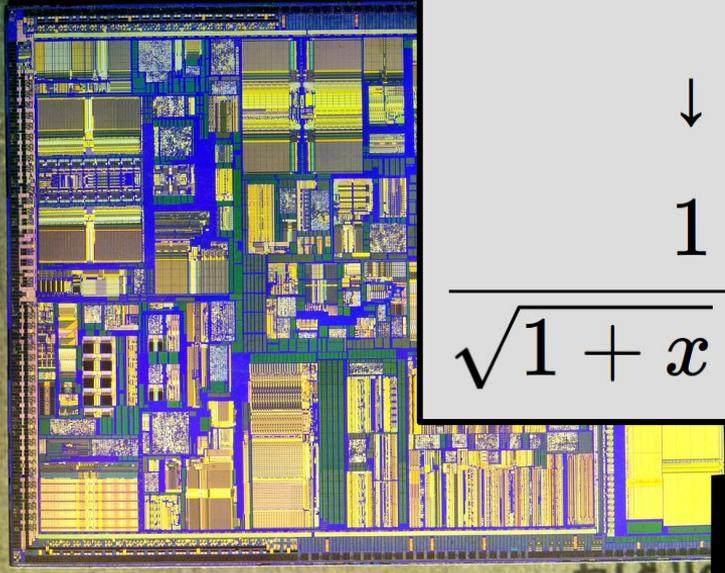
2023

PL tools

```
17 string sInput;
18 int iLength, iN;
19 double dblTemp;
20 bool again = true;
21
22 while (again) {
23     iN = -1;
24     again = false;
25     getline(cin, sInput);
26     system("cls");
27     stringstream(sInput) >> dblTemp;
28     iLength = sInput.length();
29     if (iLength < 4) {
30         again = true;
31         continue;
32     } else if (sInput[iLength - 3] != '.') {
33         again = true;
34         continue;
35     } while (++iN < iLength) {
36         if (isdigit(sInput[iN])) {
37             continue;
38         } else if (iN == (iLength - 3)) {
39             continue;
40         }
41     }
42     // ... (rest of the code)
43 }
```

next-gen PL tools

```
17 string sInput;  
18 int iLength, iN;  
19 double dblTemp;  
20 bool again = true;  
21  
22 while (again) {  
23     iN = -1;  
24     again = false;  
25     getline(cin, sInput);  
26     system("cls");  
27     stringstream(sInput) >> dblTemp;  
28     iLength = sInput.length();  
29     if (iL  
30         ag  
31         co  
32     } else  
33         ag  
34         cc  
35         whil  
36         if
```

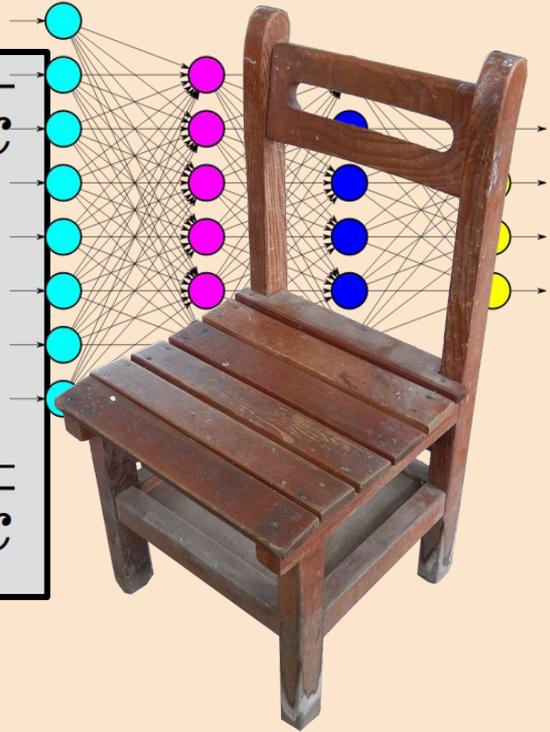


$$\sqrt{x + 1} - \sqrt{x}$$

↓

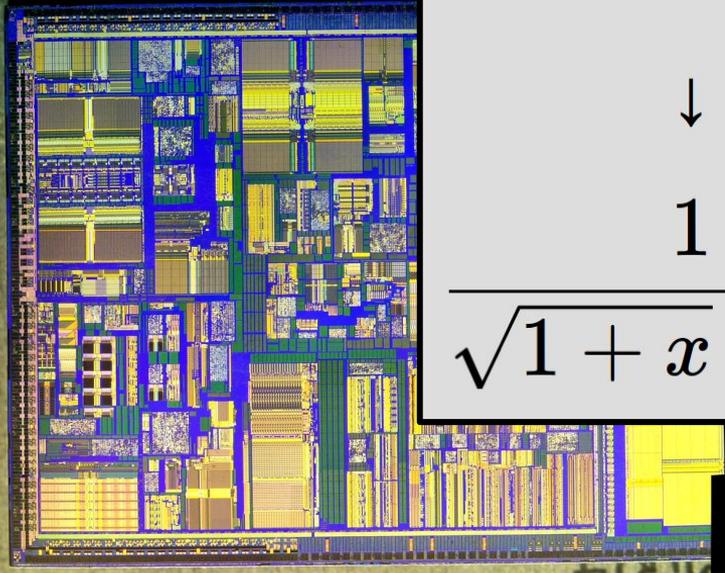
$$1$$

$$\sqrt{1 + x} + \sqrt{x}$$

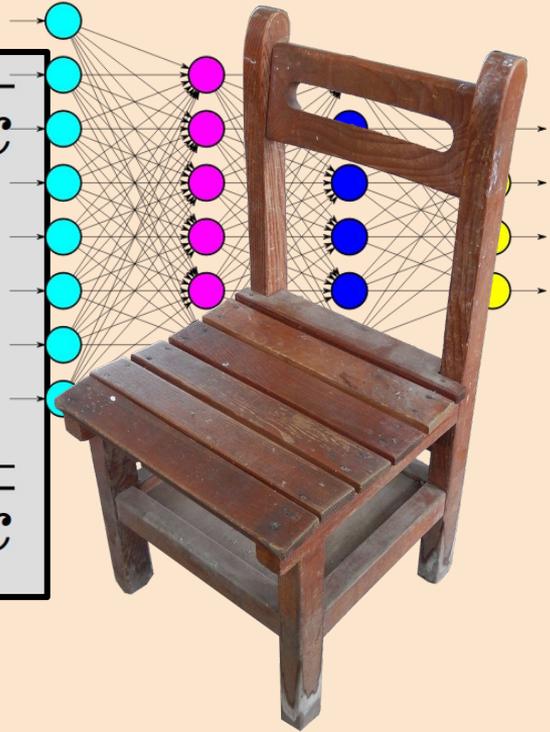


PL is everywhere

```
17 string sInput;  
18 int iLength, iN;  
19 double dblTemp;  
20 bool again = true;  
21  
22 while (again) {  
23     iN = -1;  
24     again = false;  
25     getline(cin, sInput);  
26     system("cls");  
27     stringstream(sInput) >> dblTemp;  
28     iLength = sInput.length();  
29     if (iL  
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32     } else  
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35         whil  
36         if
```



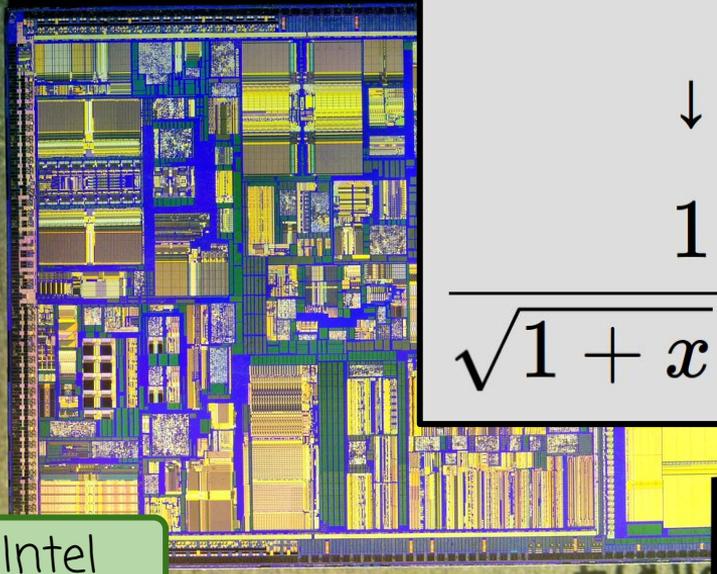
$$\frac{\sqrt{x+1} - \sqrt{x}}{1} = \frac{1}{\sqrt{1+x} + \sqrt{x}}$$



equational reasoning

Fastly

```
18 int iLength;  
19 double dblTemp;  
20 bool again = true;  
21  
22 while (again) {  
23     in = -1;  
24     again = false;  
25     getline(cin, sInput);  
26     system("cls");  
27     stringstream(sInput) >> dblTemp;  
28     iLength = sInput.length();  
29     if (iL  
30         ag  
31         co  
32     } else  
33         ag  
34         co  
35     } whil  
36     if
```



Intel

Herbie

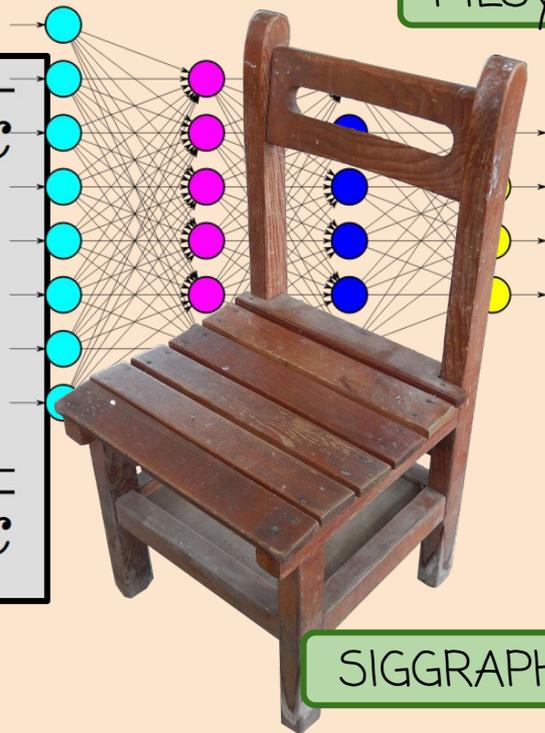
$$\sqrt{x + 1} - \sqrt{x}$$

↓

$$1$$

$$\sqrt{1 + x} + \sqrt{x}$$

MLSys '21



SIGGRAPH '22

$$(a * 2) / 2$$



a

$$(a * 2) / 2$$

=

a

$$(a * 2) / 2 \rightarrow a$$

rewrite it!

useful

$$(x * y) / z = x * (y / z)$$

$$x / x = 1$$

$$x * 1 = x$$

not so useful

$$x * 2 = x \ll 1$$

$$x * y = y * x$$

$$x = x * 1$$

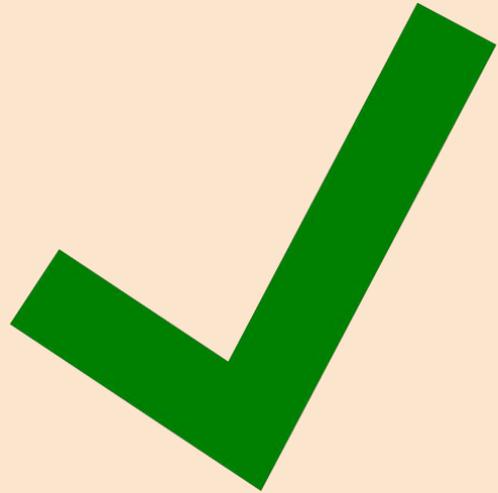
$$(a * 2) / 2 \rightarrow a * (2 / 2) \rightarrow a * 1 \rightarrow a$$

happy path

$$(x * y) / z = x * (y / z)$$

$$x / x = 1$$

$$x * 1 = x$$



$(a * 2) / 2 \Rightarrow (a \ll 1) / 2$ ❌ wrong turn

$(a * 2) / 2 \Rightarrow (2 * a) / 2 \Rightarrow (a * 2) / 2$ loop

$a \Rightarrow a * 1 \Rightarrow a * 1 * 1 \Rightarrow \dots$ infinite size

pitfalls

$$x * 2 = x \ll 1$$

$$x * y = y * x$$

$$x = x * 1$$

but critical for
other inputs

$(a * 2) / 2 \rightarrow (a \ll 1) / 2$ **✗** wrong turn

$(a * 2) / 2 \rightarrow (2 * a) / 2 \rightarrow (a * 2) / 2$ **loop**

Whitfield & Soffa 1997

“the order of applying code transformations [...] can have an impact on the quality of code produced”

but critical for other inputs

$$x * 2 = x \ll 1$$

$$x * y = y * x$$

$$x = x * 1$$

$$(a * 2) / 2 \rightarrow a$$

which rewrite? when?

useful

$$(x * y) / z = x * (y / z)$$

$$x / x = 1$$

$$x * 1 = x$$

not so useful

$$x * 2 = x \ll 1$$

$$x * y = y * x$$

$$x = x * 1$$

$$x / x = 1$$



$$x * 2 = x \ll 1$$

$$x = x * 1$$



$$x = \ll * x$$

which rewrite? when?

all of them! all the time!



$$(z | k) * x = z | (k * x)$$



$$x * y = y * x$$

$x / x = 1$

$x * 2 = x \ll 1$

$x = x * 1$

$x = 1 * x$

e-graphs
+
equality saturation

al

!

$(z | h) * x = z | (h * x)$

$x * y = y * x$

talk in a slide

⊕ background : e-graphs and eqsat

- egg: fast, flexible eqsat

POPL 21★

- intermission : community & applications

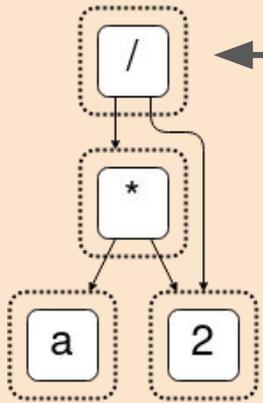
PLDI 20, MLSys 21, OOPLSA 22★, SIGGRAPH 22, FMCAD 22, POPL 23

- current work : connections to database

POPL 22

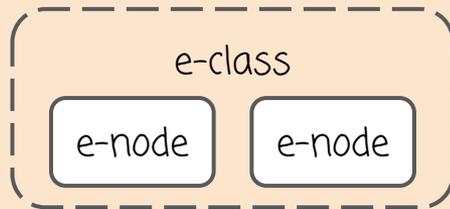
- big picture

e-graphs?



← this e-class represents

$$(a * 2) / 2$$



e-graphs?

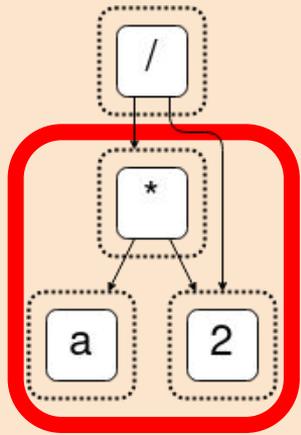
Knuth '76, Kozen '77, Downey '80, ... , Nelson '80

- context-free grammars, finitely presented algebras, and/or graphs \rightarrow e-graphs
- designed for use inside theorem provers

e-node

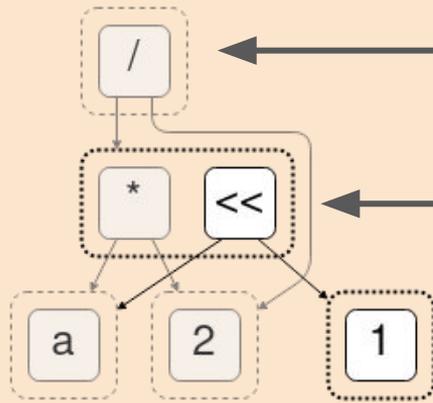
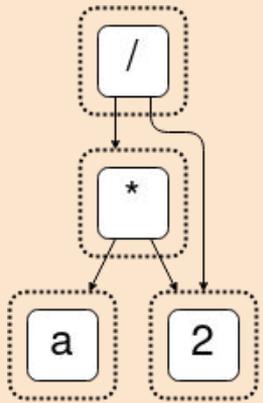
e-node

growing an e-graph



$$x * 2 \rightarrow x \ll 1$$

growing an e-graph

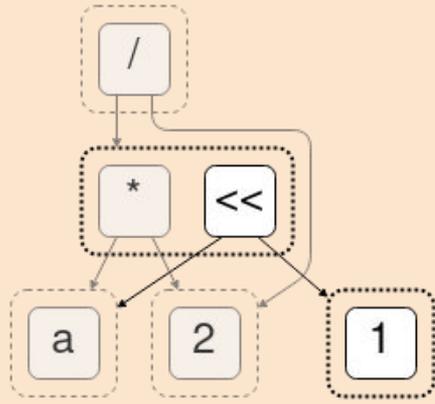
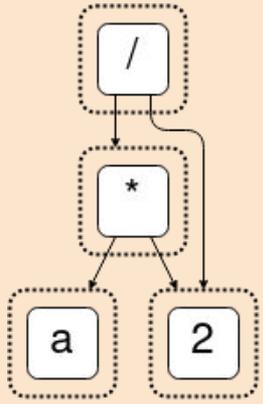


this e-class represents
 $(a * 2) / 2$ and $(a << 1) / 2$

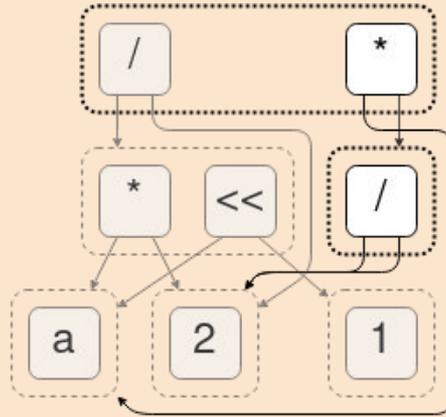
this e-class represents
 $(a * 2)$ and $(a << 1)$

$$x * 2 \rightarrow x << 1$$

growing an e-graph

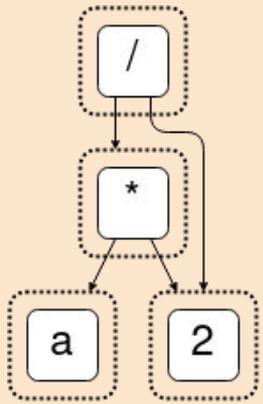


$x * 2 \rightarrow x \ll 1$

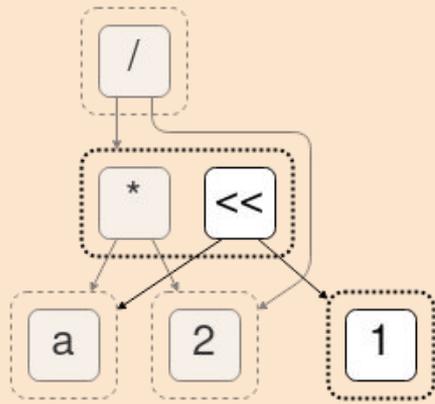


$(x * y) / z \rightarrow x * (y / z)$

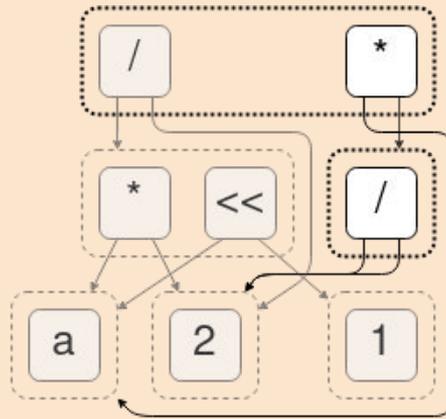
e-graphs are compact



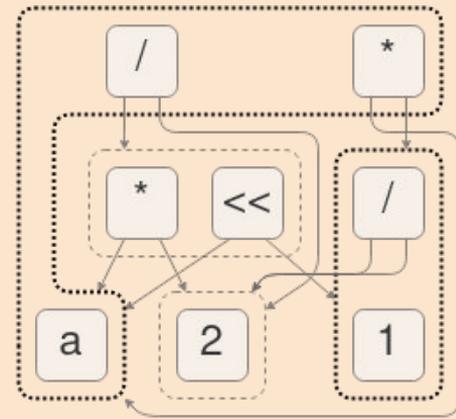
$$x * 2 \rightarrow x \ll 1$$



$$(x * y) / z \rightarrow x * (y / z)$$



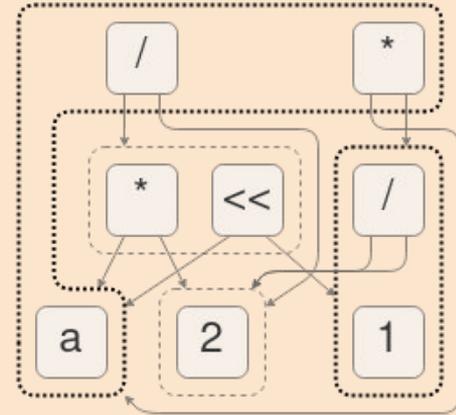
$a, a * 1,$
 $a * 1 * 1, \dots$



$$x / x \rightarrow 1$$
$$x * 1 \rightarrow x$$

saturation

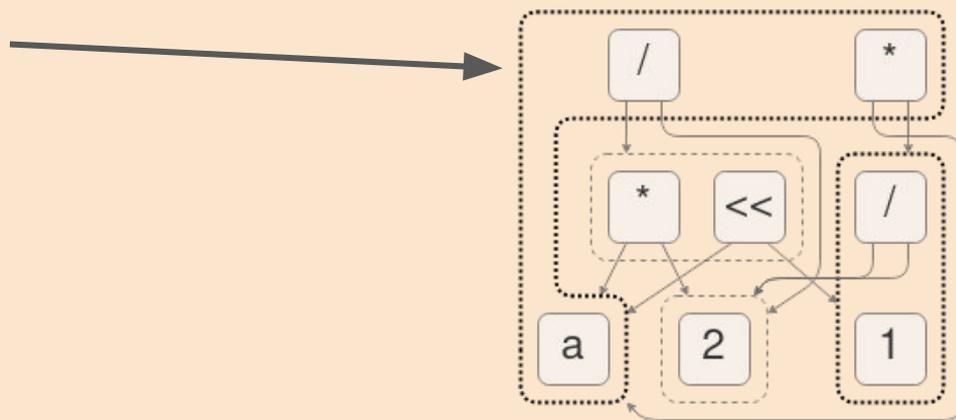
- ✓ $x * 2 \rightarrow x \ll 1$
- ✓ $(x * y) / z \rightarrow x * (y / z)$
- ✓ $x / x \rightarrow 1$
- ✓ $x * 1 \rightarrow x$



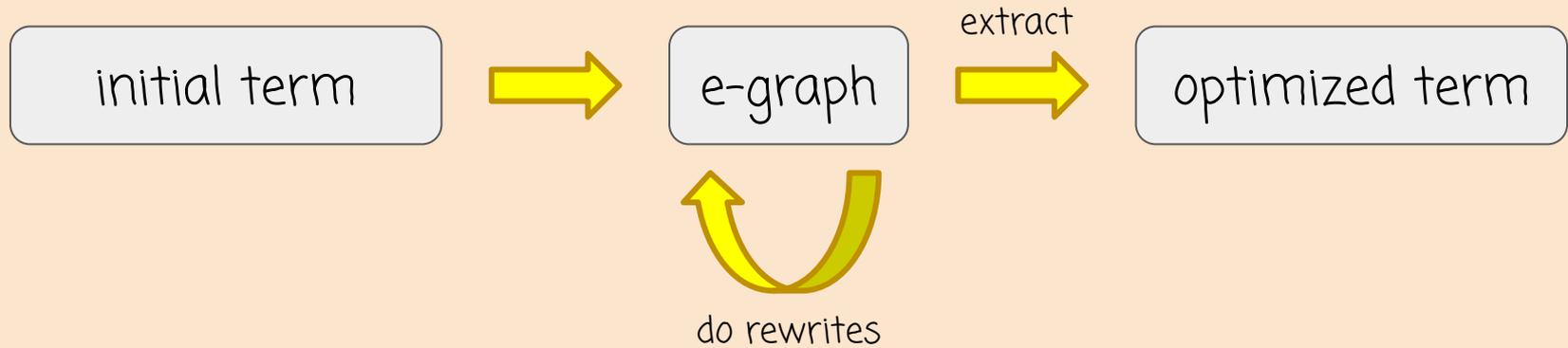
extraction

this e-class represents
 $(a * 2) / 2$, a , $a * 1$, ...

pick the smallest
(cheapest) one



equality saturation



equality saturation

EqSat [Tate et al. 2009]

- PEG encoding: algebraic IR for loops
 - Introduced "theta" nodes for sequences
- specialized to Java/LLVM, ad-hoc optimizations

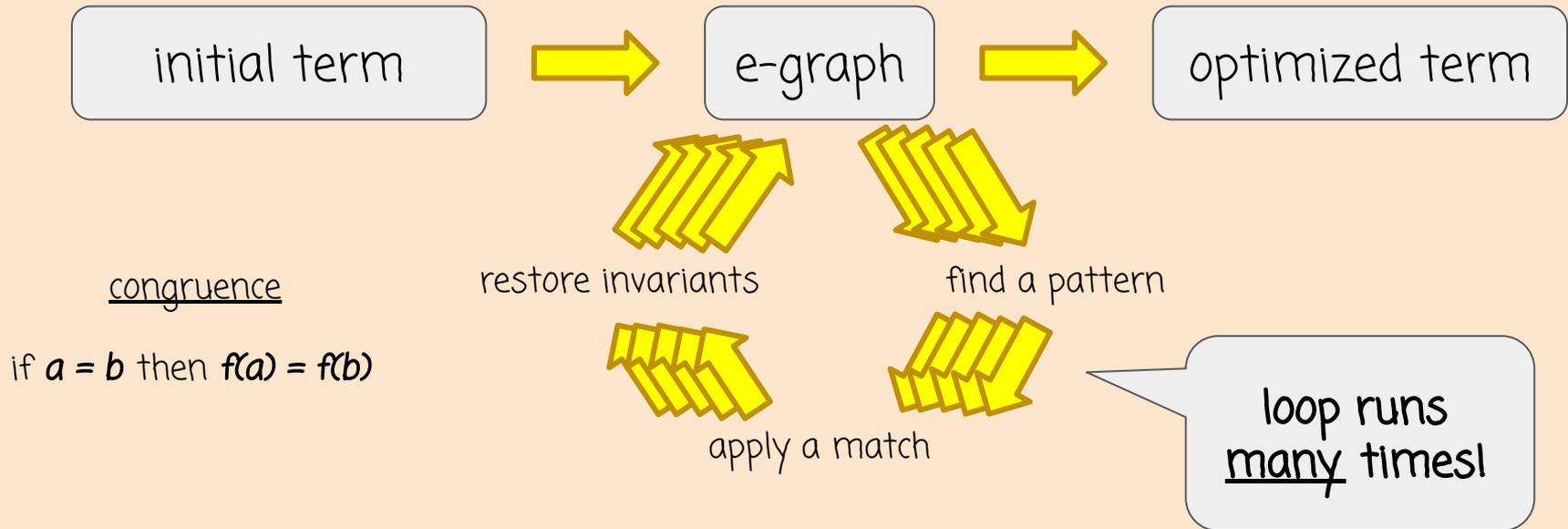
talk in a slide

- background : e-graphs and eqsat
-  egg: fast, flexible eqsat

POPL 21★

- intermission : community & applications
- current work : connections to databases
- going forward

equality saturation



equality saturation

```
def equality_saturation(expr, rewrites):  
    egraph = initial_egrgraph(expr)
```

```
    while not egraph.is_saturated_or_timeout():  
        for rw in rewrites:
```

```
            for (subst, ec) in egraph.ematch(rw.lhs):
```

```
                ec2 = egraph.add(rw.rhs.subst(subst))  
                egraph.merge(ec, ec2)
```

```
    return egraph.extract_best()
```

read

write

restore invariant

egg's equality saturation

```
def equality_saturation(expr, rewrites):  
    egraph = initial_egraph(expr)
```

```
    while not egraph.is_saturated():  
        for rw in rewrites:
```

```
            for (subst, ec) in egraph.ematch(rw.lhs):  
                ec2 = egraph.add(rw.rhs.subst(subst))  
                egraph.merge(ec, ec2)
```

easy to
parallelize

invariants broken

```
    return egraph.extract_best()
```

once per iteration

```
def equality_saturation(expr, rewrites):  
    egraph = initial_egraph(expr)
```

```
    while not egraph.is_saturated_or_timeout():  
        matches = []
```

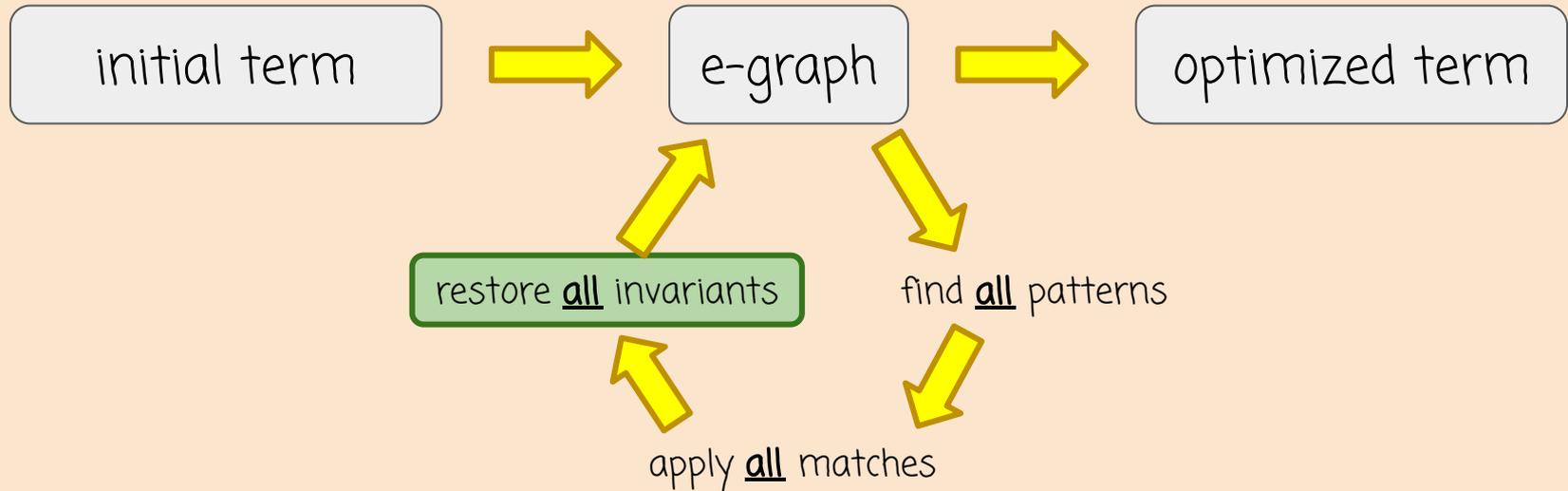
```
        for rw in rewrites:  
            for (subst, ec) in egraph.ematch(rw.lhs):  
                matches.append((rw, subst, ec))
```

```
        for (rw, subst, ec) in matches:  
            ec2 = egraph.add(rw.rhs.subst(subst))  
            egraph.merge(ec, ec2)
```

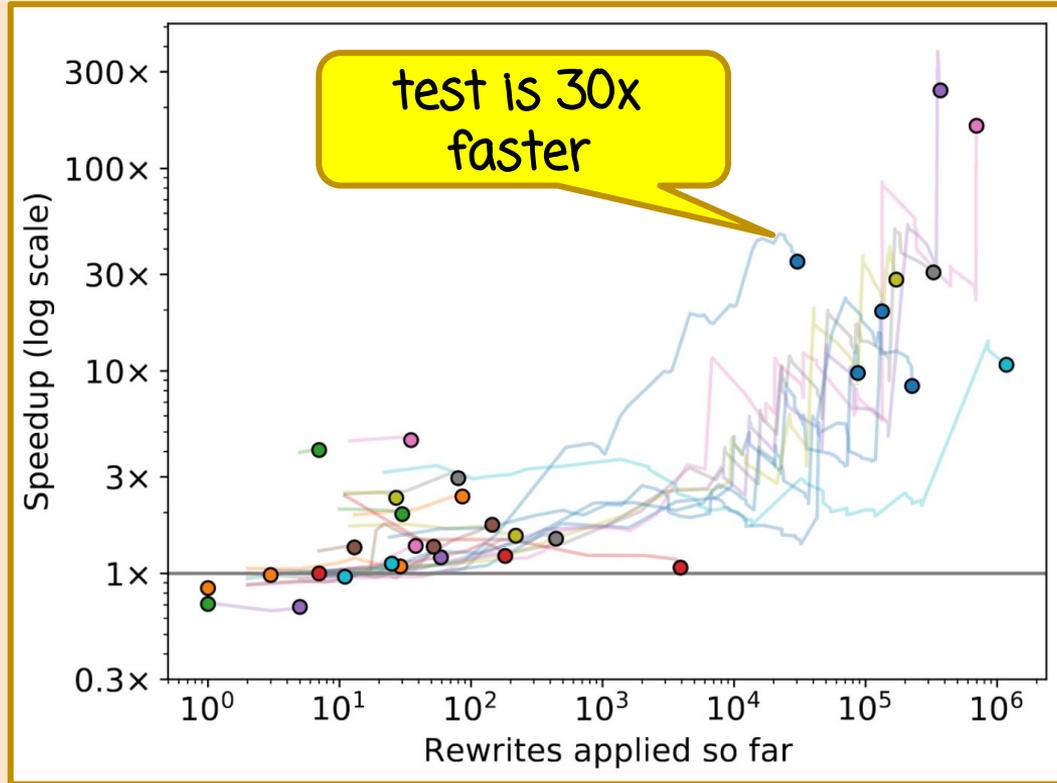
```
        egraph.rebuild()
```

```
    return egraph.extract_best()
```

deferred invariant maintenance



rebuilding is faster



fast and flexible

- there's more than syntactic rewriting
- sometimes, it's useful to consider semantics
 - $17 + 32 \rightarrow 49, \dots$
- constant folding, nullability, tensor shape, non-zero, interval arithmetic, etc, ...

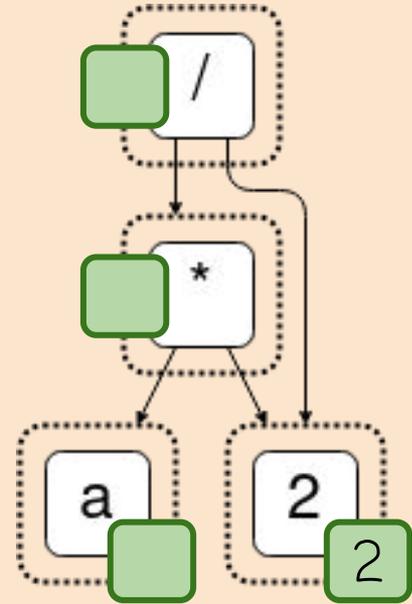
fast and flexible

analyses modulo equality

- uniform interface that works in many cases
- an understanding of analyses mean

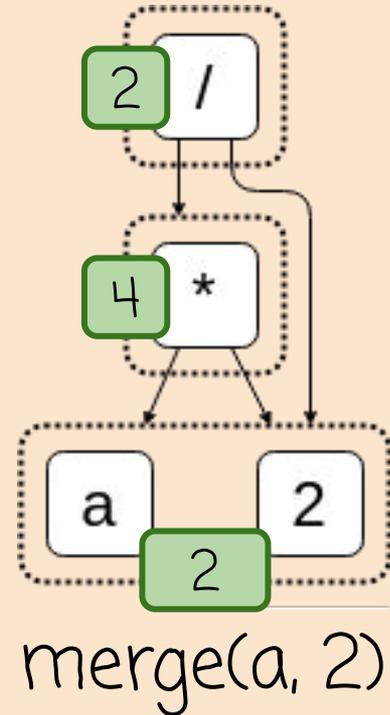
constant folding

- Option<Number> per eclass
- try to eval new e-nodes
- Option "or" on merge



constant folding

- Option<Number> per eclass
- try to eval new e-nodes
- Option "or" on merge
- it propagates up!

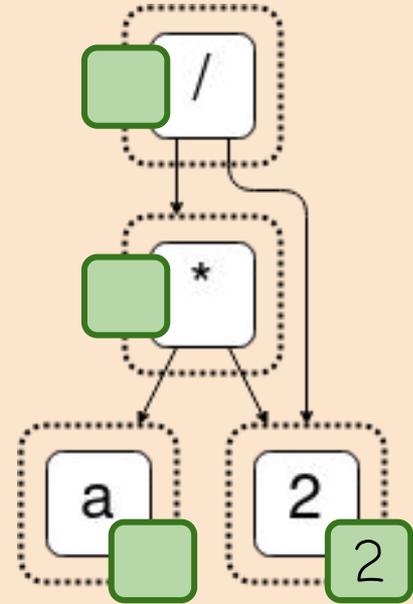


e-class analysis

- 1 fact per e-class from a join-semilattice D
- $\text{make}(n) \rightarrow d_c$
 - make a new analysis value for a new e-node
- $\text{join}(d_{c_1}, d_{c_2}) \rightarrow d_c$
 - combine two analysis values
- $\text{modify}(c) \rightarrow c'$
 - change the e-class (optionally)

constant folding

- $D = \text{Option}\langle\text{Number}\rangle$
- `make = eval`
- `join = option "or"`
- `modify = add the constant`



e-class analysis uses

- lift program analyses to e-graphs
- conditional & dynamic rewrites
 - $x / x = 1$ iff $x \neq 0$
- can express other e-graph "hacks"
 - debugging, pruning, on-the-fly extraction

e-class analysis invariant

for each e-class

fixed point

$$\forall c \in G. \quad d_c = \bigvee_{n \in c} \text{make}(n) \quad \text{and} \quad \text{modify}(c) = c$$

Analysis data is
LUB
(lattice properties)

egg: fast & easy e-graphs

- Rust library for generic e-graphs and eqsat
- packaged and documented: <https://docs.rs/egg>
- tutorials, industrial and academic users

talk in a slide

- background : e-graphs and eqsat
- egg: fast, flexible eqsat
- ⊕ intermission : community & applications

PLDI 20, MLSys 21, OOPLSA 22★, SIGGRAPH 22, FMCAD 22, POPL 23

- current work : connections to databases
- big picture

some egg projects

- ARITH 2021 Herbie: floating point 3000x faster
- VLDB 2020 SPORES: linear algebra 1.2-5x better
- MLSys 2021 Tensat: ML compute graphs 23% better, 48x faster
- SIGMOD 2022 Infer loop invs. for query opt up to 10,000x faster
- OOPSLA 2021* Ruler: Theory synthesis better rulesets, 25x faster
- ARITH 2022 Arithmetic datapath opt up to 71% smaller
- Intel, AWS, Fastly, Certora, ...

detour: intervals

x in $[0, 1]$
 y in $[1, 2]$

$x + y$ in $[1, 3]$

detour: intervals

x in $[0, 1]$
 y in $[1, 2]$

$$1 - 2y / (x + y) \quad \text{in} \quad [-3, 1/3]$$

$$= (x - y) / (x + y) \quad \text{in} \quad [-2, 0]$$

$$= 2x / (x + y) - 1 \quad \text{in} \quad [-1, 1]$$

intervals modulo equality

$$\begin{array}{l} x \text{ in } [0, 1] \\ y \text{ in } [1, 2] \end{array}$$

$$\begin{array}{l} 1 - 2y / (x + y) \text{ in } [-3, 1/3] \\ = (x - y) / (x + y) \text{ in } [-2, 0] \\ = 2x / (x + y) - 1 \text{ in } [-1, 1] \end{array} \left. \vphantom{\begin{array}{l} 1 - 2y / (x + y) \\ = (x - y) / (x + y) \\ = 2x / (x + y) - 1 \end{array}} \right\} [-1, 0]$$

intervals modulo equality

x in
y in

Expression	Initial	Improved	Width Change
$x^2 - 2x + 1$	$[-2, 3]$	$[0, 1]$	-80%
$x(2 - xy) - \frac{1}{y}$	$[-5, 1.5]$	$[-4.5, 0]$	-31%
$\sqrt{x+1} - \sqrt{x}$	$[0, \sqrt{2}]$	$\left[\frac{1}{\sqrt{2}+\sqrt{3}}, \frac{1}{1+\sqrt{2}}\right]$	-93%
$\frac{x}{x+y}$	$\left[\frac{1}{4}, 1\right]$	$\left[\frac{1}{4}, \frac{3}{4}\right]$	-33%

1/3]

0]

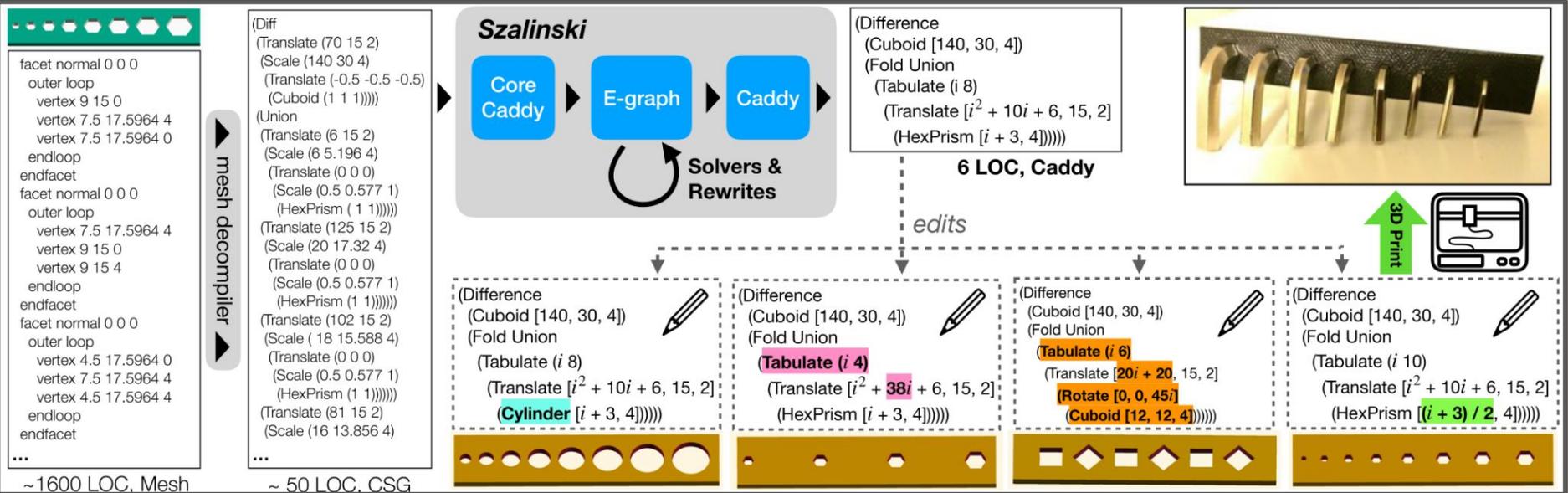
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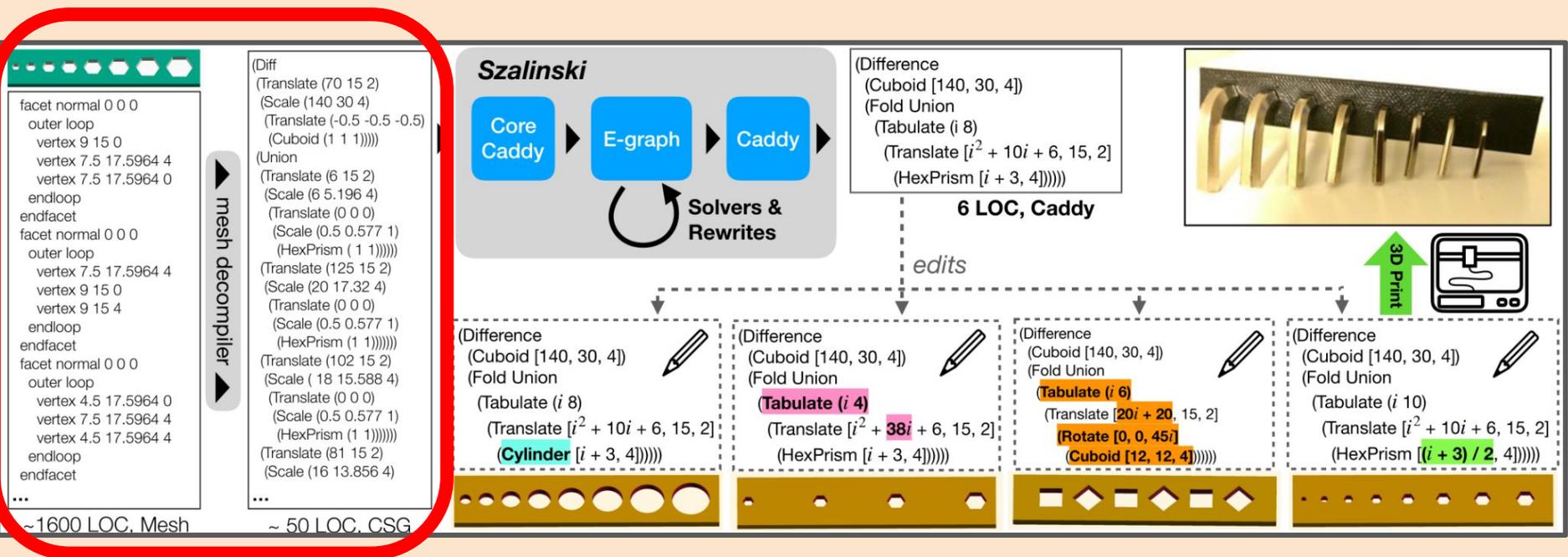
Sam Coward et. al, EGRAPHS 22



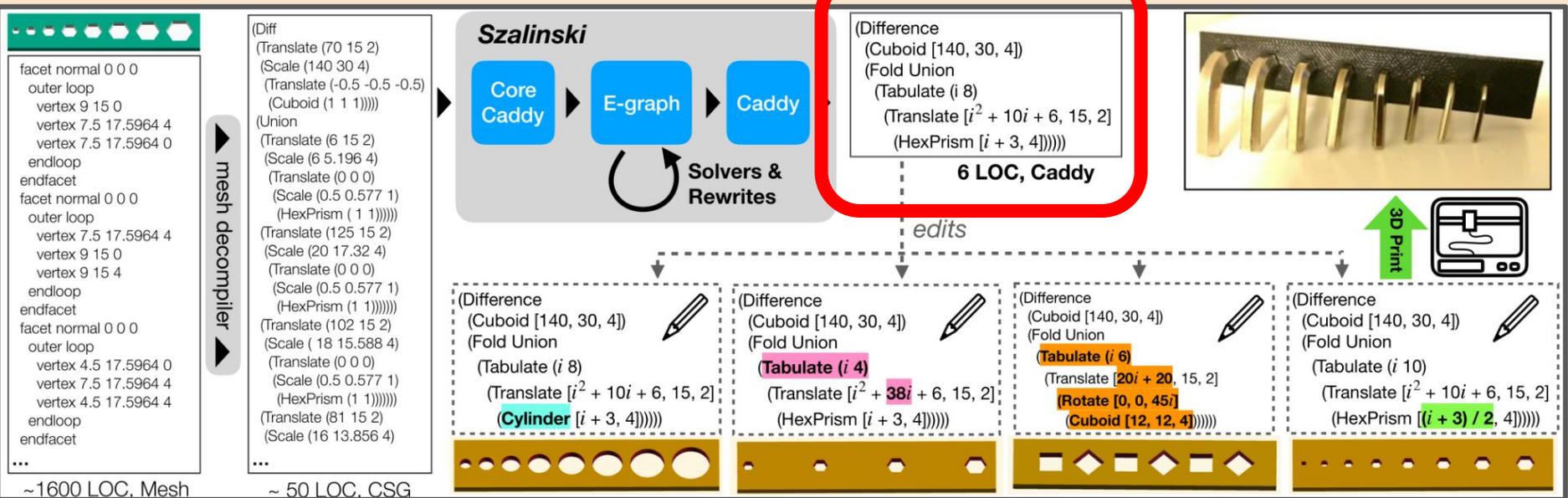
Szalinski



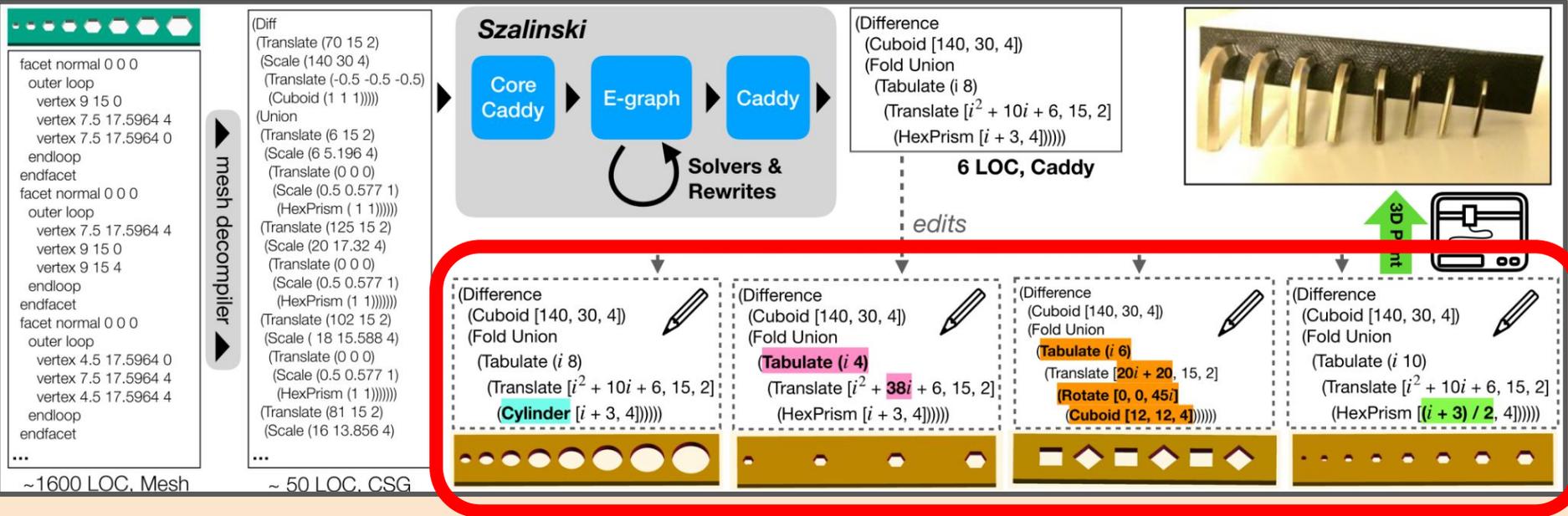
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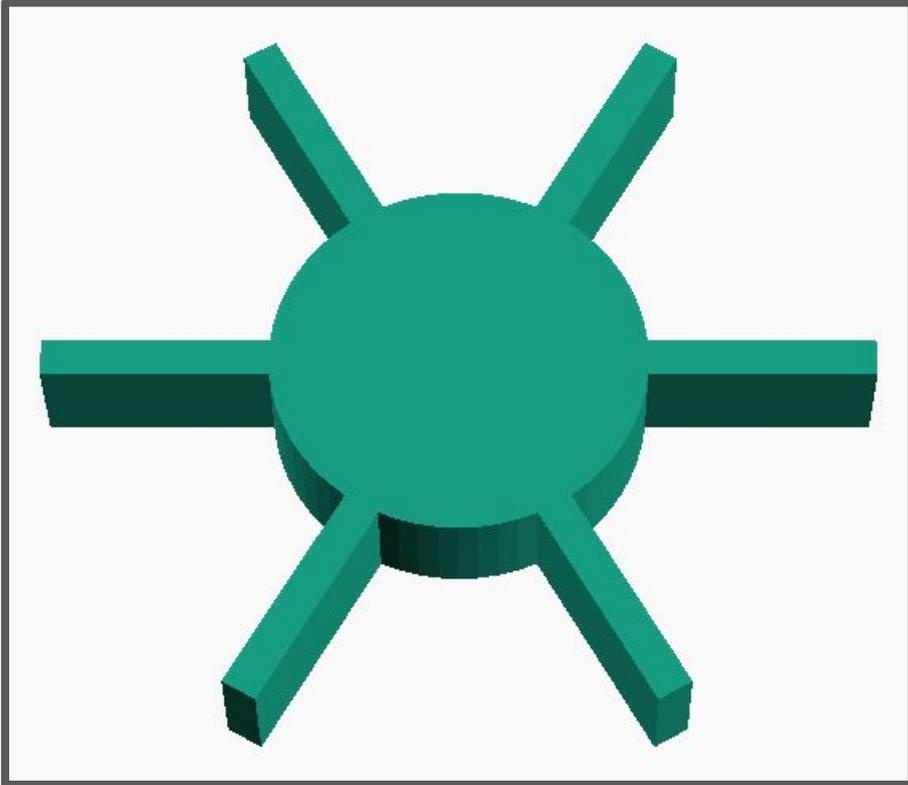
Szalinski



Szalinski

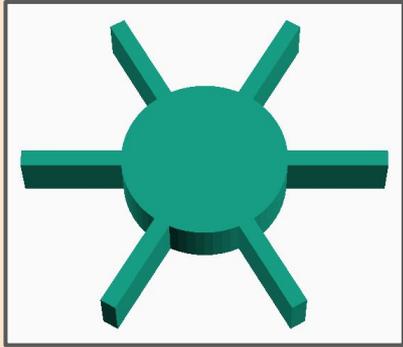


what we want



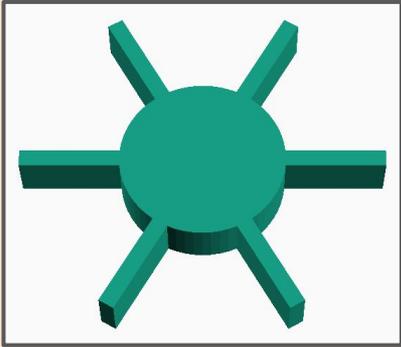
```
(Union  
  (Cylinder [1, 5, 5])  
  (Fold Union  
    (Tabulate (i 6)  
      (Rotate [0, 0, 60 * i]  
        (Translate [1, -0.5, 0]  
          (Cuboid [10, 1, 1]))))))))
```

easy mode



```
(Union  
  (Cylinder [1,5])  
  (Union  
    (Rotate [0,0,0] (Translate [1,-0.5,0] (Cuboid [10,1,1])))  
    (Rotate [0,0,60] (Translate [1,-0.5,0] (Cuboid [10,1,1])))  
    (Rotate [0,0,120] (Translate [1,-0.5,0] (Cuboid [10,1,1])))  
    (Rotate [0,0,180] (Translate [1,-0.5,0] (Cuboid [10,1,1])))  
    (Rotate [0,0,240] (Translate [1,-0.5,0] (Cuboid [10,1,1])))  
    (Rotate [0,0,300] (Translate [1,-0.5,0] (Cuboid [10,1,1])))
```

hard mode



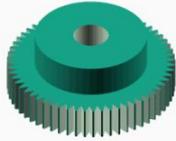
(Union

```
(Rotate [0,0,120] (Translate [1,-0.5,0] (Cuboid [10,1,1])))  
(Scale [10,1,1] (Translate [0.1,-0.5,1] (Cuboid [1,1,1])))  
(Rotate [0,0,300] (Translate [1,-0.5,0] (Cuboid [10,1,1])))  
(Scale [5,5,1] (Cylinder [1,1]))  
(Translate [-1,0.5,0] (Scale [-1,-1,1] (Cuboid [10,1,1])))  
(Rotate [0,0,240] (Translate [1,-0.5,0] (Cuboid [10,1,1])))  
(Rotate [0,0,60] (Translate [1,-0.5,0] (Cuboid [10,1,1])))
```

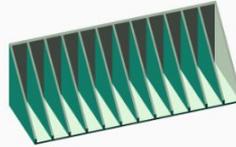
$$x = (\text{Rotate } [0 \ 0 \ 0] \ x)$$

Szalinski

PLDI 20



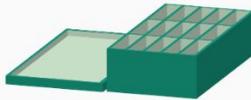
```
(Fold Difference
(List (Union
  (Cylinder [100, 80, 80])
  (Cylinder [50, 120, 120]))
(Translate [0, 0, -1] (Cylinder [102, 25, 25]))
(Fold Union (Tabulate (i 60)
  (Rotate [0, 0, 6 * i]
  (Translate [125, 0, 0]
  (Scale [2.5, 1, 1]
  (Rotate [0, 0, 45]
  (Translate [0, 0, 25]
  (Cuboid [10, 10, 52]))))))))))))
```



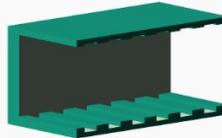
```
(Fold Union
(Tabulate (i 12)
(Translate [0, 13 * i, 0]
(Fold Difference
(List
(Cuboid [53.1 14.5 58])
(Translate [1.5, 1.5, 1.5]
(Cuboid [51.6, 11.5, 56.6]))
(Translate [0 0 58]
(Rotate [0, 45, 0]
(Cuboid [101.5, 14.5, 100]))))))))
```



```
(Fold Union
(Tabulate (i 10) (j 5)
(Translate
[12.2 * i + 12.2, 12.2 * j + 12.2, 0]
(Difference
(Cylinder [13, 7.1, 7.1])
(Translate [0, 0, 3]
(Cylinder [11, 5.1, 5.1]))))))))
```



```
(Union (Difference
(Cuboid [60, 120, 30])
(Fold Union
(Tabulate (i 5) (j 3)
(Translate [12 * i + 2, 39.3 * j + 2, 2]
(Cuboid [9.6 37.3 28]))))
(Fold Difference
(Map2 Translate
(List [-67, -2, 0] [-65, 0, 2])
(List (Cuboid [65, 125, 6])
(Cuboid [60, 120, 4]))))))))
```



```
(Difference
(Cuboid [57, 30, 30])
(Difference
(Translate [0, 5, 1.5] (Cuboid [57, 25, 27]))
(Fold Union
(Map2 Translate
(Tabulate (i 7) (j 2) [9 * i, 5, 28 - 26.5 * j]
(Concat
(List (Tabulate (i 6) (j 2)
(Cuboid [4.5, 25, j + 0.5]))
(List (Cuboid [3, 25, 0.5])
(Cuboid [3, 25, 1.5]))))))))
```

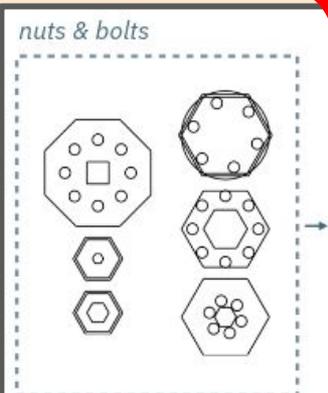


```
(Difference
(Translate [-57, -3, 3]
(Cuboid [117, 75, 175]))
(Fold Union
(Tabulate (i 10)
(Translate [-51, -3, 16 * i + 6]
(Cuboid [105, 58, 13]))))))
```

Babble

POPL 23

nuts & bolts



```
ngon = λsize sides →  
  T (repeat (T 1 (M 1 θ -0.5 (0.5 / tan (π / sides)))) sides  
    (M 1 ((2 * π) / sides) θ θ))  
  (M size θ θ θ)
```

scaled n-gon

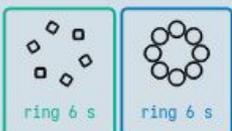


```
con_hex = λinner_size →  
  C (ngon 4 6) (ngon inner_size 6)
```



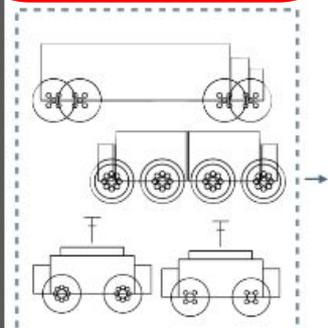
concentric scaled hexagons

```
ring = λn shape →  
  repeat (offset 1.5 shape) n (rotate n)
```

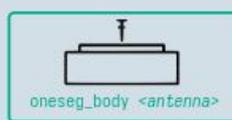


ring of shapes

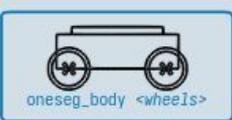
vehicles



```
oneseg_body = λrest →  
  C (C (C (C (T (T (r_s θ θ) (xform_x θ)) (xform_x -8))  
    (T (T (r_s 16 4.5) (M 1 θ θ 2.25))  
    (xform_x θ)))  
    (T (T (r_s θ θ) (xform_x θ)) (xform_x 8)))  
    (T (r_s 12 1) (M 1 θ θ 5))) rest
```



oneseg_body <antenna>



oneseg_body <wheels>

one-segment vehicle body

```
circle_ring = λn → repeat  
  (T (T c (M 0.25 θ θ θ))  
    (M 1 θ θ 0.53033 0.53033))  
  n  
  (M 1 ((2 * π) / xθ) θ θ)
```



circle_ring 4 *circle_ring 8*

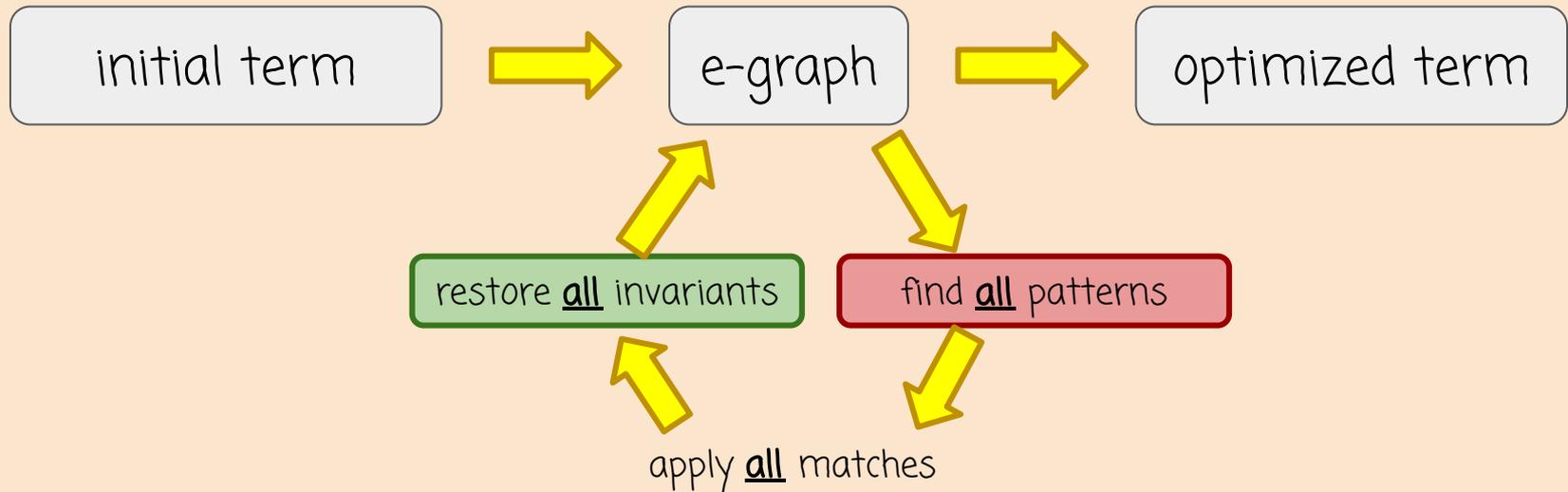
ring of circles

talk in a slide

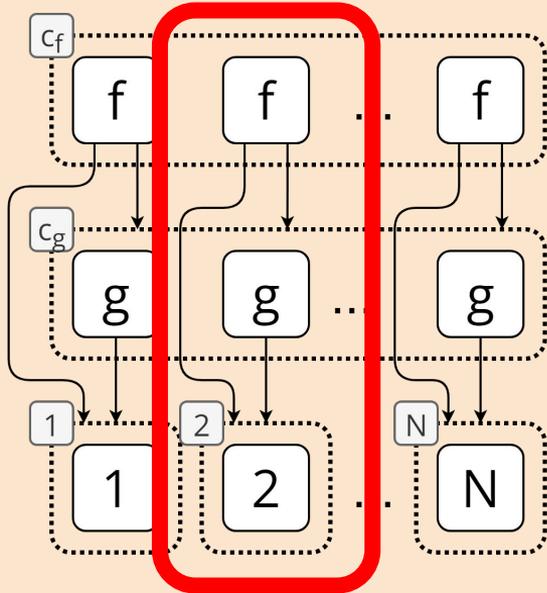
- background : e-graphs and eqsat
- egg: fast, flexible eqsat
- intermission : community & applications
- (=) current work : connections to databases
- big picture

POPL 22

equality saturation



e-matching



f g e-nodes

e-classes

pattern

$f(a, g(a))$

subst

$\{a \mapsto 1\}$

$\{a \mapsto 2\}$

...

$\{a \mapsto N\}$

terms

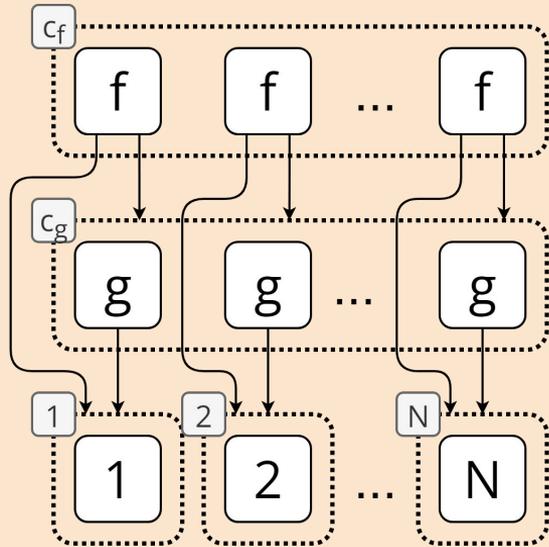
$f(1, g(1))$

$f(2, g(2))$

...

$f(N, g(N))$

e-matching



pattern

$f(\alpha, g(\alpha))$

for e-class c in e-graph E :

```
for f-node  $n_1$  in  $c$ :
```

```
  subst = {root  $\mapsto$   $c$ ,  $\alpha \mapsto n_1.child_1$ }
```

```
  for g-node  $n_2$  in  $n_1.child_2$ :
```

```
    if subst[ $\alpha$ ] =  $n_2.child_1$ :
```

```
      yield subst
```

N^2 time, but only N matches

e-matching

- existing impls are backtracking based & complex
- doesn't help with equality constraints
- no data complexity results
 - NP-hard in pattern size... e-graph size??

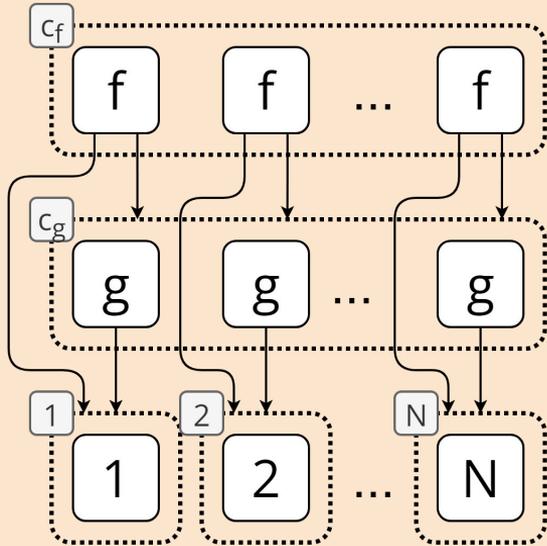
e-matching

Efficient E-matching [de Moura & Bjørner 2007]

- state-of-the art in Z3
- compile to VM, many complex optimizations
- same as egg (for now...)

○ NP-hard in pattern size... e-graph size??

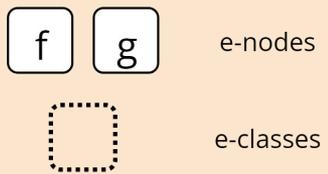
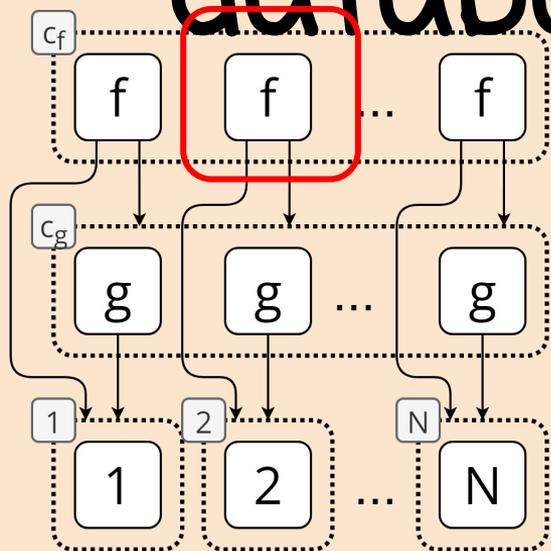
e-graph



  e-nodes

 e-classes

e-graph = database



R_f		
id	arg ₁	arg ₂
c_f	1	c_g
c_f	2	c_g
...
c_f	N	c_g

e-matching = db query

pattern

conjunctive query

$f(a, g(a))$

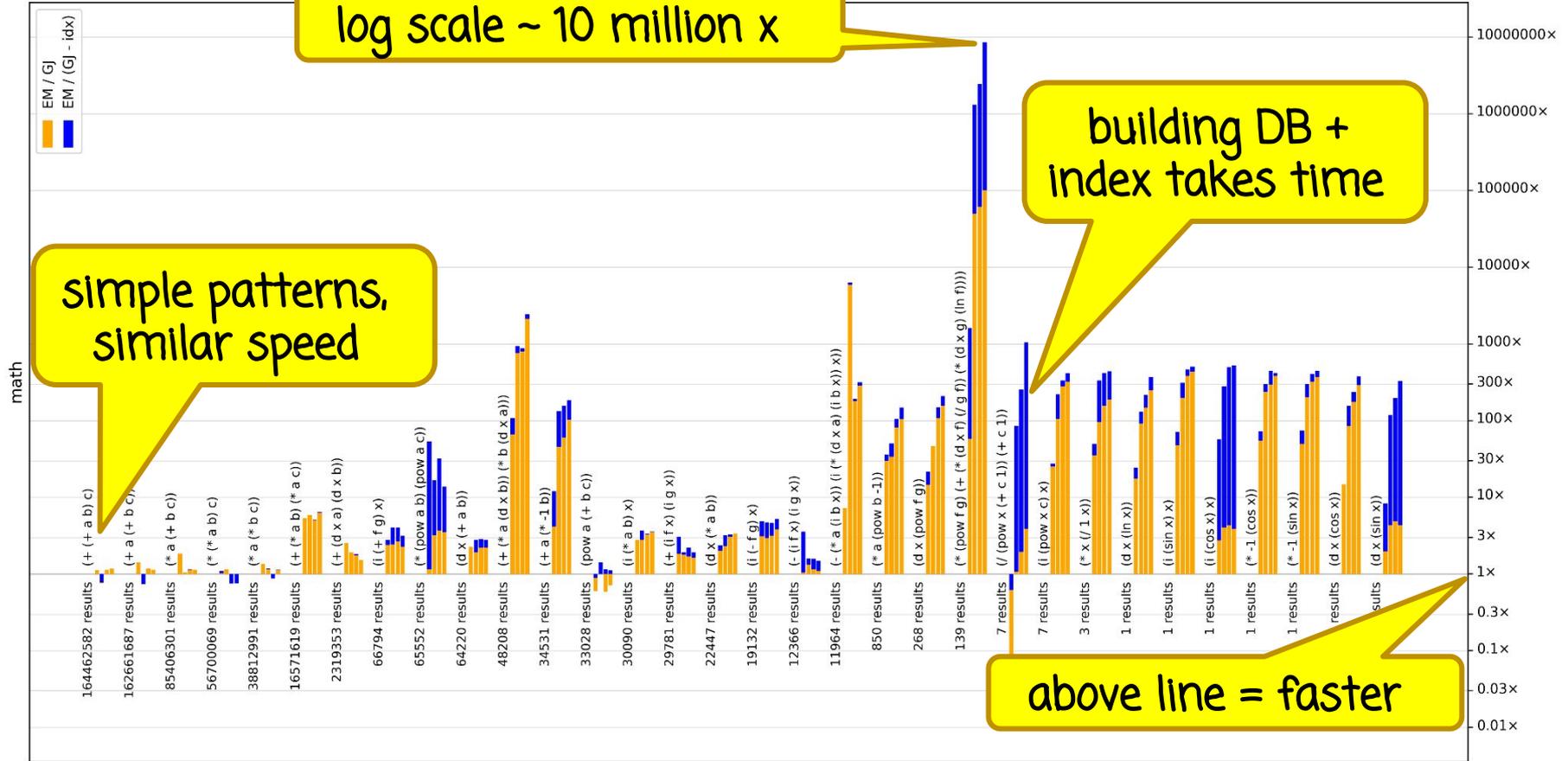
$Q(\text{root}, \alpha) \leftarrow R_f(\text{root}, \alpha, \boxed{x}), R_g(\boxed{x}, \alpha)$

log scale ~ 10 million x

building DB + index takes time

simple patterns, similar speed

above line = faster



data complexity results

THEOREM 9. *Relational e-matching is worst-case optimal; that is, fix a pattern p , let $M(p, E)$ be the set of substitutions yielded by e-matching on an e-graph E with N e-nodes, relational e-matching runs in time $O(\max_E(|M(p, E)|))$.*

w.r.t "worst-case" e-graph of size N

THEOREM 10. *Fix an e-graph E with N e-nodes that compiles to a database I , and a fix pattern p that compiles to conjunctive query $Q(\bar{X}) \leftarrow R_1(\bar{X}_1), \dots, R_m(\bar{X}_m)$. Relational e-matching p on E runs in time $O\left(\sqrt{|Q(I)| \times \prod_i |R_i|}\right) \leq O\left(\sqrt{|Q(I)| \times N^m}\right)$.*

w.r.t a specific e-graph

data complexity results

THE
the se
runs i

Worst-Case Optimal Join [Ngo, et. al, 2014 & 2018]

E) be
tching

- beautiful, recent DB result
- efficient join algo for cyclic queries
- Generic Join (go learn it!)

THE
that c
in time

tern p
E runs

w.r.t a specific e-graph

db freebies

- simple implementation
 - separate optimization pass
- multipatterns
 - $a \times b = \text{split}(1, (a ++ c) \times b)$
 - $a \times c = \text{split}(2, (a ++ c) \times b)$
- incrementality via semi-naive

talk in a slide

- background : e-graphs and eqsat
- egg: fast, flexible eqsat
- intermission : community & applications
- current work : connections to databases
- ☺ big picture

big picture

- PL tools need equational reasoning
- e-graphs let you work with multiplicity
- egg is a sweet spot: fast, flexible, usable
- make e-graphs the tool for equational reasoning

what's next?

- fundamental advances
 - e-graphs + datalog, contextual equality, ...
- big applications
 - JIT, theorem proving, query optimization
- building community
 - NSF Grant, EGRAPHS 2023, courses

thank you!

mwillsey.com

egraphs-good.github.io

EGRAPHS 2023 @ PLDI
Deadline April 5

